# A NEW PROOF OF THE $d$-CONNECTEDNESS OF $d$-POLYTOPES 

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#### Abstract

Balinski has shown that the graph of a $d$-polytope is $d$ connected. In this note we give a new proof of Balinski's theorem.


A fundamental result in the theory of convex polytopes is the following theorem of M. Balinski [1]:

Theorem. The graph $\mathcal{G}(P)$ of a d-polytope $P$ is $d$-connected.
Here a $d$-polytope is a convex polytope $P$ in $\mathbf{R}^{n}$ whose affine hull aff $P$ has dimension $d$. The graph $\mathcal{G}(P)$ is the graph whose vertices are the vertices of $P$ and whose edges are the edges of $P$. A graph $\mathcal{G}$ is $d$-connected provided that is has at least $d+1$ vertices and the removal of as many as $d-1$ vertices (and the edges incident to a removed vertex) does not destroy connectedness. For background information on convex polytopes, including Balinski's theorem, the reader may consult A. Brøndsted [2] and B. Grünbaum [3].

The proofs of Balinski's theorem in [2] and [3] are variants of Balinski's original proof. The aim of this note is to present a new proof based on a different idea.

We need a little terminology. Let $P$ be a $d$-polytope. Identifying aff $P$ with $\mathbf{R}^{d}$, we may assume that $P$ lies in $\mathbf{R}^{d}$. For any facet, i.e. $(d-1)$-dimensional face, $F$ of $P$ we denote by $K(F)$ the closed supporting halfspace of $P$ bounded by the hyperplane aff $F$. It is a standard fact that

$$
\begin{equation*}
P=\bigcap\left\{K(F) \mid F \in \mathcal{F}_{d-1}(P)\right\} \tag{1}
\end{equation*}
$$

where $\mathcal{F}_{d-1}(P)$ denotes the set of facets of $P$, see e.g. [2, Corollary 9.6].
For any vertex $v$ of $P$ we denote by $C(P, v)$ the intersection of all closed halfspaces $K(F)$ such that $F$ is a facet of $P$ and $v$ is a vertex of $F$. It is trivial that

$$
\begin{equation*}
P \subseteq C(P, v) \subseteq K(F) \tag{2}
\end{equation*}
$$

for all facets $F$ of $P$ and all vertices $v$ of $F$. Combining (1) and (2) we obtain

$$
\begin{equation*}
P=\bigcap\{C(P, v) \mid v \in \operatorname{vert} P\} \tag{3}
\end{equation*}
$$

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where vert $P$ denotes the set of vertices of $P$.
Our proof of Balinski's theorem is based on the following lemma which shows that we may omit as many as $d-1$ vertices from the right hand side of (3):

Lemma. Let $V$ be a set of at most $d-1$ vertices of a $d$-polytope $P$ in $\mathbf{R}^{d}$. Then

$$
\begin{equation*}
P=\bigcap\{C(P, v) \mid v \in(\operatorname{vert} P) \backslash V\} \tag{4}
\end{equation*}
$$

Proof. The inclusion $\subseteq$ is a trivial consequence of (3). To prove $\supseteq$, let $x$ be a point not in $P$. Then by (1) there is a facet $F$ of $P$ such that $x \notin K(F)$. Being a ( $d-1$ )-polytope, $F$ has at least $d$ vertices, whence at least one vertex $v$ of $F$ is not in $V$. Using (2) we see that $x \notin C(P, v)$. Hence, $x$ is not in the right hand side of (4). This completes the proof of the lemma.

We next turn to
Proof of the Theorem. Let $V$ be any set of at most $d-1$ vertices of $P$. We shall prove that the subgraph $\mathcal{G}^{\prime}$ of $\mathcal{G}(P)$ spanned by (vert $\left.P\right) \backslash V$ is connected. Let $\mathcal{G}^{\prime \prime}$ be a connected component of $\mathcal{G}^{\prime}$; we reach the desired conclusion by showing that $\mathcal{G}^{\prime}=\mathcal{G}^{\prime \prime}$.

We denote the vertex set of a graph $\mathcal{G}$ by vert $\mathcal{G}$. Then vert $\mathcal{G}^{\prime}=($ vert $P) \backslash V$. Let $Q$ denote the convex polytope spanned by (vert $\left.\mathcal{G}^{\prime \prime}\right) \cup V$; then vert $Q=\left(\right.$ vert $\left.\mathcal{G}^{\prime \prime}\right) \cup V$.

Let $v \in$ vert $\mathcal{G}^{\prime \prime}=($ vert $Q) \backslash V$. Since $\mathcal{G}^{\prime \prime}$ is a connected component of $\mathcal{G}^{\prime}$, any vertex of $P$ adjacent to $v$ is in (vert $\mathcal{G}^{\prime \prime}$ ) $\cup V=$ vert $Q$. In other words:
(5) Any $v \in($ vert $Q) \backslash V$ has the same adjacent vertices in $Q$ as in $P$.

The following is a standard fact, cf. [2, Corollary 11.7]:
For any convex polytope $R$ and any vertex $u$ of $R$, the affine hull of $u$ and its adjacent vertices in $R$ is aff $R$.

Now, since $P$ is a $d$-polytope, it follows from (5) and (6) that $Q$ is also a $d$-polytope. Noting that $V \subseteq$ vert $Q$, we may then apply the lemma to $Q$, obtaining

$$
\begin{equation*}
Q=\bigcap\{C(Q, v) \mid v \in(\operatorname{vert} Q) \backslash V\} \tag{7}
\end{equation*}
$$

We next claim that

$$
\begin{equation*}
P \subseteq C(Q, v) \text { for all } v \in(\text { vert } Q) \backslash V \tag{8}
\end{equation*}
$$

To prove (8), let $v \in$ (vert $Q) \backslash V$. Let $x$ be an arbitrary point of $P$. If $x=v$, then $x \in C(Q, v)$ as desired. If $x \neq v$, we choose a hyperplane $H$ which separates $v$ from the remaining vertices of $P$ and from $x$. The sets $P^{\prime}:=H \cap P$ and $Q^{\prime}:=H \cap Q$ are then $(d-1)$-polytopes, cf. [2, Theorem 11.2]. The vertices of $P^{\prime}$ are the points where the edges of $P$ connecting $v$ and its adjacent vertices in $P$ intersect $H$, cf. [2, Theorem
11.2]. Similarly for $Q^{\prime}$. It then follows from (5) that $P^{\prime}$ and $Q^{\prime}$ have the same vertices, whence $P^{\prime}=Q^{\prime}$. Let $x^{\prime}$ be the point where the segment from $v$ to $x$ intersects $H$. Then obviously $x^{\prime} \in P^{\prime}$. Since $P^{\prime}=Q^{\prime}$, it follows that $x^{\prime} \in Q^{\prime}$, whence $x^{\prime} \in Q$, and so $x^{\prime} \in C(Q, v)$. Since $C(Q, v)$ is a cone with vertex $v$, it follows that the entire halfline emanating from $v$ and passing through $x^{\prime}$ is in $C(Q, v)$. In particular, $x \in C(Q, v)$ as desired.

Combining (7) and (8) we obtain $P \subseteq Q$, whence

$$
\begin{equation*}
P=Q \tag{9}
\end{equation*}
$$

since $Q \subseteq P$ is obvious.
Finally, let $v \in$ vert $\mathcal{G}^{\prime}=($ vert $P) \backslash V$. Then $v \in($ vert $Q) \backslash V=$ vert $\mathcal{G}^{\prime \prime}$ by (9). Hence vert $\mathcal{G}^{\prime}=\operatorname{vert} \mathcal{G}^{\prime \prime}$. By the nature of $\mathcal{G}^{\prime}$ and $\mathcal{G}^{\prime \prime}$ this implies $\mathcal{G}^{\prime}=\mathcal{G}^{\prime \prime}$ as desired.

## References

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