

COUNTABLE J_a^S -ADMISSIBLE ORDINALS

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§ 0. Introduction

In [3], Platek constructs a hierarchy of jumps J_a^S indexed by elements a of a set \mathcal{O}^S of ordinal notations. He asserts that a real $X \subseteq \omega$ is recursive in the superjump S if and only if it is recursive in some J_a^S . Unfortunately, his assertion is not correct as is shown in [1]. In [1], it also has been shown that an ordinal $> \omega$ is J_a^S -admissible if it is $|a|_S$ -recursively inaccessible, where $|a|_S$ is the ordinal denoted by a .

Let A be an arbitrary set. We say that an ordinal α is A -admissible if the structure $\langle L_\alpha[A], \in, A \cap L_\alpha[A] \rangle$, which we will denote by $L_\alpha[A]$ for simplicity, is admissible, a model of the Kripke-Platek set theory KP , where $L_\alpha[A]$ is the sets constructible relative to A in fewer than α steps. We use ω_1^A or $\omega_1(A)$ to denote the first A -admissible ordinal $> \omega$, and use $\omega_1(A_1, \dots, A_n)$ for $\omega_1(\langle A_1, \dots, A_n \rangle)$.

The purpose of this paper is to prove the following theorem.

THEOREM 1. *Suppose $a \in \mathcal{O}^S$ and $\alpha > \omega$ is a countable $|a|_S$ -recursively inaccessible ordinal. Then, there exists a real $X \subseteq \omega$ such that $\alpha = \omega_1(J_a^S, X)$.*

In the case $|a|_S = 0$, $J_a^S = {}^2E$, the Kleene object of type 2, and $\omega_1({}^2E, X) = \omega_1^X$ for all reals $X \subseteq \omega$. α is an admissible ordinal if and only if it is 0-recursively inaccessible. Therefore, Theorem 1 is an extension of the following theorem of Sacks.

THEOREM 2 (Sacks [4]). *If $\alpha > \omega$ is a countable admissible ordinal, then there exists a real X such that $\alpha = \omega_1^X$.*

Sacks also showed that the real X mentioned in Theorem 2 can be taken to have the minimality property:

$$\omega_1^Y < \alpha \text{ for every } Y \text{ of lower hyperdegree than } X.$$

Received July 8, 1983.

Likewise, we can show that for every countable $|a|_s$ -recursively inaccessible $\alpha > \omega$ there is a real X such that:

$$\alpha = \omega_1(J_a^s, X);$$

and

$$\omega_1(J_a^s, Y) < \alpha \text{ for every } Y \text{ of lower } J_a^s\text{-degree than } X.^1)$$

Theorem 1 will be proved by the forcing with J_a^s -pointed perfect trees. Let $\alpha > \omega$ be a countable $|a|_s$ -recursively inaccessible ordinal and X be a generic real with respect to this forcing relation. Then $L_a[X]$ is admissible and $\alpha \leq \omega_1(J_a^s, X)$. To see $\omega_1(J_a^s, X) \leq \alpha$, we must show that X preserves sufficiently many admissible ordinals below α to make α to be $\langle J_a^s, X \rangle$ -admissible.

§1. $|a|_s$ -recursively inaccessible ordinals

A normal type 2 object is a total function F from ω^ω to ω such that the Kleene object 2E of type 2:

$${}^2E(f) = \begin{cases} 0 & \text{if } (\exists n)f(n) = 0, \\ 1 & \text{otherwise,} \end{cases}$$

is recursive in F . The superjump $S(F)$ of F is a type 2 object defined by:

$$S(F)(\langle n, f \rangle) = \begin{cases} 0 & \text{if } \{n\}^F(f) \text{ is defined,} \\ 1 & \text{otherwise.} \end{cases}$$

Platek [3] defines a hierarchy J_a^s of type 2 objects along with a set \mathcal{O}^s of ordinal notations, starting from 2E and iterating the superjump operation transfinitely.

An ordinal α is 0-recursively inaccessible if it is admissible. α is $(\sigma+1)$ -recursively inaccessible if it is σ -recursively inaccessible and a limit of σ -recursively inaccessible ordinals. For limit λ , α is said to be λ -recursively inaccessible if it is σ -recursively inaccessible for all $\sigma < \lambda$. Let X be an arbitrary set. σ -recursively-in- X inaccessible ordinals are defined in the same way starting from X -admissible ordinals. By $RI(\sigma, X)$, we denote the class of all σ -recursively-in- X inaccessible ordinals. In the case $X = \emptyset$, $RI(\sigma, \emptyset)$ is the class of all σ -recursively inaccessible ordinals.

The following lemma, due to Aczel and Hinman, gives a characterization of $\omega_1(J_a^s, X)$ for $X \subseteq \omega$.

1) For J_a^s -degrees, the reader may refer to [5].

LEMMA 3 (Aczel and Hinman [1]). *Suppose $a \in \mathcal{O}^S$ and $\sigma = |a|_S$, the ordinal denoted by a . Then $\sigma < \omega_1(J_a^S)$, and for any ordinal $\alpha > \omega$ and $X \subseteq \omega$:*

$$\alpha \in RI(\sigma, X) \rightarrow \alpha \text{ is } \langle J_a^S, X \rangle\text{-admissible,}$$

and $\omega_1(J_a^S, X)$ is the least ordinal in $RI(\sigma, X)$.

Let λ_0 be the least ordinal λ such that λ is λ -recursively inaccessible. Lemma 3 shows that $|\mathcal{O}^S| = \sup \{|a|_S : a \in \mathcal{O}^S\} \leq \lambda_0$. In [1], it has been shown that $|\mathcal{O}^S| = \lambda_0$.

Let $\alpha > \omega$ be a countable admissible ordinal. Using the unbounded Levy forcing over L_α , we can add to L_α a generic function $K: (\alpha - \omega) \times \omega \rightarrow \alpha$ such that if $\omega \leq \beta < \alpha$ then the function $\lambda nK(\beta, n)$ is a bijection from ω onto β . Therefore, in $L_\alpha[K]$ all sets are countable. It has been shown in [4] that $\langle L_\alpha[K], \in, K \rangle$ is an admissible structure in which Σ_1 -DC (Σ_1 -Dependent Choice) holds.

Suppose $a \in \mathcal{O}^S$. For any $X, Y \subseteq \omega$, $X \leq_{J_a^S} Y$ means X is recursive in $\langle J_a^S, Y \rangle$, which is equivalent to that $X \in L_\rho[J_a^S, Y]$, where $\rho = \omega_1(J_a^S, Y)$. X and Y have the same J_a^S -degree, $X \equiv_{J_a^S} Y$, if $X \leq_{J_a^S} Y$ and $Y \leq_{J_a^S} X$. $X <_{J_a^S} Y$ if $X \leq_{J_a^S} Y$ but $X \not\equiv_{J_a^S} Y$.

LEMMA 4. *Suppose $\alpha > \omega$ is a countable $|a|_S$ -recursively inaccessible ordinal and K is a generic function with respect to the unbounded Levy forcing over L_α . Then for any $X, Y \subseteq \omega$:*

$$X \leq_{J_a^S} Y \text{ and } Y \in L_\alpha[K] \longrightarrow X \in L_\alpha[K].$$

Proof. The unbounded Levy forcing preserves admissible ordinals. That is, if $\beta < \alpha$ is an admissible ordinal then β is K -admissible. This is because for admissible β , $K \upharpoonright (\beta - \omega) \times \omega$ is generic with respect to the unbounded Levy forcing over L_β . Therefore, if $Y \in L_\alpha[K]$ then α is $|a|_S$ -recursively-in- Y inaccessible, so $L_\rho[Y] \subseteq L_\alpha[K]$, where $\rho = \omega_1(J_a^S, Y)$. Thus we have the lemma. □

§ 2. J_a^S -pointed perfect trees

Let a be an element of \mathcal{O}^S such that $|a|_S > 0$. We put $J = J_a^S$ for simplicity.

A perfect tree is a set P of finite sequences of 0's and 1's such that:

$$(1) \quad p \in P \text{ and } q \subseteq p \longrightarrow q \in P;$$

and

$$(2) \quad (\forall p \in P)(\exists q, r \in P) (q \text{ and } r \text{ are incompatible extensions of } p),$$

where $q \subseteq p$ denotes that p is an extension of q . For a perfect tree P , $[P]$ denotes the set of all infinite paths through P :

$$[P] = \{f \in 2^\omega : (\forall n) \bar{f}(n) \in P\}.$$

We say that P is J -pointed if:

$$(3) \quad (\forall f \in [P])(\omega_1(J, P) \leq \omega_1(J, f) \text{ and } P \in L_{\omega_1(J, P)}[f]).$$

Note that if P is J -pointed then it is \leq_J -pointed in the sense of Sacks [4:2.1], but not vice versa.

LEMMA 5. *Suppose P is J -pointed. If $X \subseteq \omega$ and $P \leq_J X$, then there exists a J -pointed $Q \subseteq P$ such that $Q \equiv_J X$.*

Proof. In [4:2.3], Sacks constructed a perfect subtree Q of P such that:

$$(4) \quad Q \text{ is recursive in } P \text{ and } f \text{ for every } f \in [Q];$$

and

$$(5) \quad Q \equiv_J X.$$

To see Q is J -pointed in our sense, fix $f \in [Q]$. Since P is J -pointed and $f \in [P]$, by (3), we have:

$$(6) \quad P \in L_{\omega_1(J, P)}[f].$$

Clearly:

$$(7) \quad f \in L_{\omega_1(J, P)}[f].$$

From (4), (6) and (7), we obtain:

$$(8) \quad Q \in L_{\omega_1(J, P)}[f].$$

From (5) and the assumption $P \leq_J X$, we see:

$$(9) \quad \omega_1(J, P) \leq \omega_1(J, Q).$$

From (8) and (9), we obtain $Q \in L_{\omega_1(J, Q)}[f]$. □

For any ordinal δ , $\{\delta\}^J$ denotes the δ -th element^J of $L[f]$ in the canonical

wellordering on $L[f]$. A perfect tree P is said to be uniformly J -pointed if there exists an ordinal δ such that:

$$(10) \quad (\forall f \in [P])(P = \{\delta\}^f \text{ and } \delta < \omega_1(J, f)).$$

Obviously, uniformly J -pointed perfect trees are J -pointed. Let $\alpha > \omega$ be a countable $|a|_S$ -recursively inaccessible ordinal and K a generic function over L_α in the sense of the unbounded Levy forcing. Observe that if P is uniformly J -pointed and $P \in L_\alpha[K]$ then there exists a $\delta < \alpha$ which satisfies (10) since the leftmost path f_P through P is recursive in P and so $\omega_1(J, f_P) \leq \omega_1(J, P) < \alpha$.

Let M be a countable admissible set and P be a perfect tree in M . Then P becomes a partially ordered set as usual. The forcing with P as the set of conditions is called the local Cohen forcing over M and denoted by $\| \frac{P}{M}$, or simply by $\| \frac{P}{M}$. If $f \in [P]$ is generic with respect to $\| \frac{P}{M}$, then $M[f]$ is an admissible set, and so is $L_\mu[f]$, where $\mu = M \cap On$.

LEMMA 6. For any $\xi < \alpha$ and any J -pointed perfect tree P in $L_\alpha[K]$, there exists a uniformly J -pointed perfect tree $Q \subseteq P$ such that $\xi < \omega_1(J, Q)$ and $Q \in L_\alpha[K]$.

Proof. Since ξ is countable in $L_\alpha[K]$, there is a real $X \in L_\alpha[K]$ such that ξ is recursive in X . By Lemma 5, there is a J -pointed perfect subtree P_1 of P such that $P_1 \equiv_J X$. Then we see $\xi < \omega_1(J, P_1)$, and $P_1 \in L_\alpha[K]$ by Lemma 4. Thus, we may assume $\xi < \omega_1(J, P)$ from the beginning. Put $M = L_{\omega_1(J, P)}[P]$. Consider the local Cohen forcing relation $\| \frac{P}{M}$ over M . Since P is J -pointed, we have:

$$(11) \quad (\forall f \in [P]) \omega_1(J, P) \leq \omega_1(J, f);$$

and

$$(12) \quad (\forall f \in [P]) (\exists \gamma < \omega_1(J, P)) \{\gamma\}^f = P.$$

By (12), there exists a $p_0 \in P$ and $\gamma < \omega_1(J, P)$ such that:

$$(13) \quad p_0 \| \frac{P}{M} \{\check{\gamma}\}^\mathcal{F} = \check{P},$$

where \mathcal{F} is the canonical name which denotes the generic reals. As in [4: 2.10], we can construct a perfect tree $Q \subseteq P$ such that:

$$(14) \quad Q \in L_{\omega_1(J, P)}[P];$$

and

$$(15) \quad (\forall f \in [Q]) \{\dot{\gamma}\}^f = P.$$

From (14), we can find a $\delta < \omega_1(\mathcal{J}, P)$ such that $\{\delta\}^P = Q$. So, by (15), there is an $\varepsilon < \omega_1(\mathcal{J}, P)$ such that:

$$(16) \quad (\forall f \in [Q]) \{\varepsilon\}^f = Q.$$

Let f_Q be the leftmost branch of Q . Then, by (11):

$$(17) \quad \omega_1(\mathcal{J}, P) \leq \omega_1(\mathcal{J}, f_Q) \leq \omega_1(\mathcal{J}, Q).$$

Hence, from (16), we see that Q is uniformly \mathcal{J} -pointed. By (17), we also see $\xi < \omega_1(\mathcal{J}, Q)$. Since $P \in L_\alpha[K]$, we have $\omega_1(\mathcal{J}, P) \leq \alpha$, and so $Q \in L_{\omega_1(\mathcal{J}, P)}[P] \subseteq L_\alpha[K]$. \square

Let \mathcal{L} be a first-order language. A Π_1^1 formula in \mathcal{L} is a second-order formula of the form:

$$(\forall S_1) \cdots (\forall S_m) \psi,$$

where S_1, \dots, S_m are predicate variables and ψ is a first-order formula in the expanded language $\mathcal{L} \cup \{S_1, \dots, S_m\}$.

LEMMA 7. *Suppose A is a countable admissible set such that $\omega \in A$ and $\mathcal{L} \in A$ is a first-order language. Let $\theta(x_1, \dots, x_n)$ be a Π_1^1 formula in \mathcal{L} . Then there exists a Σ_1 formula $\Phi(x_1, \dots, x_n, y)$ such that for any structure $\mathcal{M} = \langle M, \dots \rangle \in A$ for \mathcal{L} and any $a_1, \dots, a_n \in M$:*

$$A \models \Phi(a_1, \dots, a_n, \mathcal{M}) \iff \mathcal{M} \models \theta(a_1, \dots, a_n).$$

Proof. This is well-known. See, e.g., Barwise [2: IV. 3.1]. \square

Using this lemma, we obtain the following lemma.

LEMMA 8. *The set of all uniformly \mathcal{J} -pointed perfect trees in $L_\alpha[K]$ is Σ_1 over $L_\alpha[K]$.*

Proof. Put $\sigma = |a|_S$, (recall that $\mathcal{J} = \mathcal{J}_a^S$). Let P be a perfect tree in $L_\alpha[K]$ and $\delta < \alpha$. Let $\beta(P, \delta, \sigma)$ denote the least admissible ordinal $\beta < \alpha$ such that $\max(\delta, \sigma, \omega) < \beta$ and $P \in L_\beta[K]$. The function β is Σ_1 over $L_\alpha[K]$. We can easily find a Π_1^1 formula θ in the language of set theory such that for any perfect tree $P \in L_\alpha[K]$:

$$P \text{ is uniformly } \mathcal{J}\text{-pointed} \iff (\exists \delta < \alpha) L_{\beta(P, \delta, \sigma)}[K] \models \theta(P, \delta, \sigma).$$

Thus the lemma follows from Lemma 7. \square

§ 3. Forcing with uniform J_a^S -pointed perfect trees

Suppose $|a|_S > 0$ and put $J = J_a^S$. Let $\alpha > \omega$ be a countable $|a|_S$ -recursively inaccessible ordinal and K a generic function with respect to the unbounded Levy forcing over L_α , which we fix throughout this section.

Let $\mathcal{L}(\alpha, \mathcal{T})$ be a ramified language containing names for all members of $L_\alpha[f]$. $\mathcal{L}(\alpha, \mathcal{T})$ includes: a numeral \bar{n} for each $n \in \omega$, unranked variables x, y, z, \dots ; ranked variables $x^\beta, y^\beta, z^\beta, \dots$ for each $\beta < \alpha$; and abstraction operator $\hat{}$. It is intended that \mathcal{T} denotes $\{n \in \omega : f(n) = 1\}$, that x ranges over $L_\alpha[f]$, that x^β ranges over $L_\beta[f]$, and that $\hat{x}^\beta\phi(x^\beta)$ denotes the set:

$$\{x \in L_\beta[f] : L_\beta[f] \models \phi(x)\}.$$

$C(\beta)$ is the set of names for elements of $L_\beta[f]$ and $C = \bigcup_{\beta < \alpha} C(\beta)$.

Let \mathcal{P} denote the set of all uniformly J -pointed perfect trees in $L_\alpha[K]$. P, Q, R, \dots denote the members of \mathcal{P} . For a ranked sentence ϕ of $\mathcal{L}(\alpha, \mathcal{T})$ and $P \in \mathcal{P}$, let $\rho(P, \phi)$ be the least admissible ordinal $\rho < \alpha$ such that $P \in L_\rho[K]$ and $\text{rank}(\phi) < \rho$. The function ρ is Σ_1 over $L_\alpha[K]$. The forcing relation $P \Vdash \phi$, where ϕ is a sentence of $\mathcal{L}(\alpha, \mathcal{T})$, is defined inductively:

- (1) ϕ is ranked. $P \Vdash \phi$ iff $(\forall f \in [P])L_{\rho(P, \phi)}[f] \models \phi$;
- (2) $\phi \vee \psi$ is not ranked. $P \Vdash \phi \vee \psi$ iff $P \Vdash \phi$ or $P \Vdash \psi$;
- (3) $(\exists x^\beta)\phi(x^\beta)$ is not ranked. $P \Vdash (\exists x^\beta)\phi(x^\beta)$ if $P \Vdash \phi(c)$ for some $c \in C(\beta)$;
- (4) $P \Vdash (\exists x)\phi(x)$ iff $P \Vdash \phi(c)$ for some $c \in C$;
- (5) ϕ is not ranked. $P \Vdash \neg \phi$ iff $(\forall Q \subseteq P) \nVdash \phi$.

Using Lemmas 7 and 8, it is easy to see that the forcing relation $P \Vdash \phi$, restricted Σ_1 sentences ϕ , is Σ_1 over $L_\alpha[K]$.

LEMMA 9. *For each P and ϕ , there exists a $Q \subseteq P$ such that $Q \Vdash \phi$ or $Q \Vdash \neg \phi$.*

Proof. In view of (5), we may assume that ϕ is ranked. By Lemma 6, we may also assume that $\phi \in L_\delta[P]$ for some P -admissible δ such that $\delta < \omega_1(J, P)$. Then, in $L_\delta[P]$, all sets are countable. Thus, in $L_\delta[P]$, we can enumerate all ranked sentences of rank $\leq \text{rank}(\phi)$:

$$\phi = \phi_0, \phi_1, \dots, \phi_n, \dots \quad (n \in \omega).$$

Let \Vdash^P be the local Cohen forcing relation over $L_\delta[P]$. In $L_\delta[P]$, we can construct a family $\langle q_s : s \in \text{Seq}(2) \rangle$ of elements of P such that:

$$(6) \quad q_s \Vdash^P \phi_n \text{ or } q_s \Vdash^P \neg \phi_n, \text{ where } n = \ell h(s);$$

and

$$(7) \quad q_{\widehat{s(0)}} \text{ and } q_{\widehat{s(1)}} \text{ are incompatible extensions of } q_s,$$

where $\text{Seq}(2)$ is the set of all finite sequences of 0's and 1's. Let $Q = \{q \in P : (\exists s) q \subseteq q_s\}$. Then by (7) Q is a perfect subtree of P . By (6), it is easy to see that $Q \Vdash \phi$ or $Q \Vdash \neg \phi$. Since $Q \in L_\delta[P]$, $Q = \{\gamma\}^P$ for some $\gamma < \delta$. Therefore Q is uniformly J -pointed because P is. \square

A real $f \in 2^\omega$ is said to be generic if for every dense subset \mathcal{D} of \mathcal{P} which is definable over $L_\alpha[K]$ there is a $P \in \mathcal{D}$ such that $f \in [P]$. For every $P \in \mathcal{P}$, there is a generic f such that $f \in [P]$. From Lemma 9, it follows that for every generic f and sentence ϕ :

$$L_\alpha[f] \models \phi \text{ iff } (\exists P)(f \in [P] \text{ and } P \Vdash \phi).$$

LEMMA 10. *If f is generic, then $L_\alpha[f]$ is admissible.*

Proof. We need to show that $L_\alpha[f]$ satisfies the \mathcal{A}_0 -Collection. Let $\phi(x, y)$ be a formula of $\mathcal{L}(\alpha, \mathcal{T})$ with no unranked quantifiers. We claim that if $P \Vdash (\forall n)(\exists y)\phi(n, y)$ then there exists a $Q \subseteq P$ and a $\beta < \alpha$ such that $Q \Vdash (\forall n)(\exists y^\beta)\phi(n, y^\beta)$. The proof of this claim is almost the same as that of [4:3.12] with some notational changes. So, we omit the proof here. From the claim, it follows that $L_\alpha[f]$ satisfies the \mathcal{A}_0 -Collection. \square

Proof of Theorem 1. Let $\alpha > \omega$ be a countable $|a|_s$ -recursively inaccessible ordinal and K be as before. Put $\sigma = |a|_s$ and $J = J_\alpha^s$. In the case $\sigma = 0$, Theorem 1 is exactly Theorem 2, which has already been established by Sacks [4]. So we may assume $\sigma > 0$. Let $f_0 \in 2^\omega$ be a generic real over $L_\alpha[K]$ with respect to the forcing with uniform J -pointed perfect trees. By Lemma 6, for each $\xi < \alpha$, the set $\{P \in \mathcal{P} : \xi < \omega_1(J, P)\}$ is dense in \mathcal{P} . It is obviously definable over $L_\alpha[K]$. Therefore there is a $P \in \mathcal{P}$ such that $f_0 \in [P]$ and $\xi < \omega_1(J, P)$. Since P is J -pointed, we have:

$$\xi < \omega_1(J, P) \leq \omega_1(J, f_0).$$

Thus, we have $\alpha \leq \omega_1(J, f_0)$. To see $\alpha = \omega_1(J, f_0)$, we must show that $\alpha \in RI(\sigma, f_0)$. At first we consider the case where $\sigma = \tau + 1$ for some τ . It

is sufficient to prove that α is a limit of ordinals in $RI(\tau, f_0)$, since then by induction on τ we can show that $\alpha \in RI(\tau, f_0)$, (note that $\alpha \in RI(0, f_0)$ by Lemma 10). Suppose $\xi < \alpha$. We shall show that the following set \mathcal{D}_ξ is dense in \mathcal{P} :

$$\mathcal{D}_\xi = \{P \in \mathcal{P} : (\exists \delta < \alpha)(\xi < \delta \text{ and } (\forall f \in [P])\delta \in RI(\tau, f))\}.$$

Assume this can be done. Using Lemma 7, it is easy to see that \mathcal{D}_ξ is Σ_1 over $L_\alpha[K]$. Therefore, for every $\xi < \alpha$, there exists a $\delta < \alpha$ such that $\xi < \delta$ and $\delta \in RI(\tau, f_0)$

To show that \mathcal{D}_ξ is dense in \mathcal{P} , take an arbitrary $P \in \mathcal{P}$. By Lemma 6, we may assume $\xi < \omega_1(J, P)$. Take a $\delta \in RI(\tau, P)$ so that $\xi < \delta < \omega_1(J, P)$. Such a δ exists because $\omega_1(J, P)$ is a limit of ordinals in $RI(\tau, P)$. Consider the local Cohen forcing relation \Vdash^P over $L_\delta[P]$. Let δ^+ be the next P -admissible ordinal of δ . Then, $L_\delta[P]$ is countable in $L_{\delta^+}[P]$. So we can enumerate inside $L_{\delta^+}[P]$ all sentences of the appropriate forcing language:

$$\phi_0, \phi_1, \dots, \phi_n, \dots \quad (n \in \omega).$$

As in the proof of Lemma 9, we can construct a perfect subtree $Q \in L_{\delta^+}[P]$ of P such that:

$$(\forall f \in [Q])f \text{ is generic with respect to } \Vdash^P.$$

Q is uniformly J -pointed since $Q \in L_{\delta^+}[P]$, $\delta^+ < \omega_1(J, P)$ and P is uniformly J -pointed. To show that $\delta \in RI(\tau, f)$ for all $f \in [Q]$, take $f \in [Q]$. Let $\beta \leq \delta$ be an arbitrary P -admissible ordinal $> \omega$, and \Vdash^β be the local Cohen forcing relation over $L_\beta[P]$. It is easy to see that f is generic with respect to \Vdash^β , and so β is f -admissible. From this, by induction on τ , we see that $\delta \in RI(\tau, f)$.

Now we consider the case where σ is a limit ordinal. The proof is carried out in the same way. For any $\xi < \alpha$ and any $\tau < \sigma$, let $\mathcal{D}_{\xi\tau}$ be the set:

$$\{P \in \mathcal{P} : (\exists \delta < \alpha)(\xi < \delta \text{ and } (\forall f \in [P])\delta \in RI(\tau, f))\}.$$

Then $\mathcal{D}_{\xi\tau}$ is dense in \mathcal{P} and definable over $L_\alpha[K]$. Therefore, we have that $\alpha = \omega_1(J, f_0)$ for any generic f_0 with respect to \Vdash . □

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