

A duality theorem for a nondifferentiable nonlinear fractional programming problem

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A duality theorem, and a converse duality theorem, are proved for a nonlinear fractional program, where the numerator of the objective function involves a concave function, not necessarily differentiable, and also the support function of a convex set, and the denominator involves a convex function, and the support function of a convex set. Various known results are deduced as special cases.

Introduction

Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$, $g : \mathbb{R}^n \rightarrow \mathbb{R}$, and $h : \mathbb{R}^n \rightarrow \mathbb{R}^m$ be continuous functions, with $-f$ and g convex. Let $S \subset \mathbb{R}^m$ be a closed convex cone, which may in particular be the nonnegative orthant \mathbb{R}_+^m ; let the function h be S -convex [6]. Let C_1 and C_2 be closed convex sets in \mathbb{R}^n . Consider the nonlinear fractional programming problem

$$(1) \quad (P): \quad \underset{x \in X_0}{\text{maximize}} \quad \frac{f(x) - s(x|C_1)}{g(x) + s(x|C_2)} \quad \text{subject to} \quad -h(x) \in S,$$

in which X_0 is an open convex set in \mathbb{R}^n , $s(\cdot|C_i)$ is the support function of the set C_i ($i = 1, 2$), and it is assumed that

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$$(2) \quad x \in X_0 \text{ (or } -h(x) \in S) \Rightarrow g(x) + s(x|C_2) > 0 .$$

Associate to (P) the problem

$$(D): \quad \text{minimize } z \text{ subject to } z \geq 0, \quad y \in S^*, \quad v \in C_1, \quad w \in C_2, \\ u \in X_0, y, z, v, w$$

$$(3) \quad 0 \in \partial(-f+zg)(u) + \partial(y^T h)(u) + (v+zw),$$

$$(4) \quad -f(u) + zg(u) + y^T h(u) + (v+zw)^T u \geq 0 .$$

In (D), S^* is the dual cone to S [6], and ∂ denotes subdifferential [15].

Under suitable hypothesis, (D) will be shown to be a dual problem to (P); under somewhat different assumptions, (P) will be shown to be a dual problem to (D).

The constraint $-h(x) \in S$ is *locally solvable* [4], [5] at x_0 if $-h(x_0) \in S$ and, for some $\delta > 0$, whenever the direction d satisfies $-h(x_0) - h'(x_0; d) \in S$ and $\|d\| < \delta$ (where $h'(x_0; d)$ denotes directional derivative in direction d), there exists a solution $x = x_0 + \alpha d + o(\alpha)$ to $-h(x) \in S$, valid for sufficiently small $\alpha > 0$.

(This requirement reduces to the Kuhn-Tucker constraint qualification for a constraint system $h_i(x) \leq 0$ ($i = 1, 2, \dots, m$)). The problem (P) will be said to satisfy a *constraint qualification* at x_0 if $-h(x_0) \in S$, and *either*

(a) Slater's constraint qualification holds, namely $-h(x) \in \text{int } S$ for some $x \in X_0$, or

(b) $-h(x) \in S$ is locally solvable at $x_0 \in X_0$, and the set

$$(5) \quad \bigcup_{s \in S^*} \{sh(x_0)\} \times \partial(sh)(x_0)$$

is closed in $\mathbb{R} \times \mathbb{R}^m$.

(The latter is automatic if S is a polyhedral cone and h is differentiable at x_0 [9].)

Duality theorem

The assumptions stated in the Introduction will be assumed throughout.

THEOREM 1. *Weak duality holds for (P) and (D), namely $\sup(P) \leq \inf(D)$. If (P) reaches a maximum at $x = x_0 \in X_0$, if $\max(P) \geq 0$, and if a constraint qualification holds for (P), then (D) reaches a minimum at some (u, z, y, v, w) with $u = x_0$, and $\max(P) = \min(D)$.*

Thus (D) is a strong dual [7], [8] to (P).

Proof. Let x be feasible for (P), and let (u, z, y, v, w) be feasible for (D). From a constraint for (D), $\theta + \psi + (v+zw) = 0$ for some $\theta \in \partial\varphi(u)$ and some $\psi \in \partial(y^T h)(u)$, where $\varphi = -f + zg$. Since $z \geq 0$ and $-f$ and g are convex, φ is convex. Then

$$\begin{aligned} & [f(x) - s(x|C_1)] - z[g(x) + s(x|C_2)] \\ &= f(x) - zg(x) - (v+zw)^T x \\ &= -\varphi(x) - (v+zw)^T x \\ &\leq -\varphi(u) - \theta^T(x-u) - (v+zw)^T(x-u) - (v+zw)^T u \text{ since } \theta \in \partial\varphi(u) \\ &= -\varphi(u) - (v+zw)^T u + \psi^T(x-u) \text{ by a constraint for (D)} \\ &\leq y^T h(u) + \psi^T(x-u) \text{ by a constraint for (D)} \\ &\leq y^T h(x) \text{ since } y^T h \text{ is a convex function} \\ &\leq 0 \text{ since } -h(x) \in S \text{ and } y \in S^*. \end{aligned}$$

By hypothesis, $g(x) + s(x|C_2) > 0$. Dividing by it,

$$[f(x) - s(x|C_1)] / [g(x) + s(x|C_2)] \leq z.$$

Hence $\sup(P) \leq \inf(D)$.

Now assume that (P) is maximized at x_0 , with $m = \max(P) \geq 0$, and also the constraint qualification. Then x_0 also maximizes

$$(6) \quad [f(x) - s(x|C_1)] - m[g(x) + s(x|C_2)]$$

subject to $-h(x) \in S$. Applying the appropriate nondifferentiable version of the Kuhn-Tucker Theorem, assuming constraint qualification (a) or (b)

(see [9], Theorem 4),

$$0 \in \partial(-f+mg)(x_0) + v + w + \partial(y^T h)(x_0) , \quad y^T h(x_0) = 0 ,$$

holds for some $y \in S^*$, $v \in \partial s(x_0|C_1)$, and $w \in \partial s(x_0|C_2)$. Then [15], $v \in C_1$, $v^T x_0 = s(x_0|C_1)$, $w \in C_2$, $w^T x_0 = s(x_0|C_2)$; therefore all constraints of (D) except (4) hold for $u = x_0$, $z = m$, y, v, w , and (4) holds also from $[f(x_0)-s(x_0|C_1)]/[g(x_0)+s(x_0|C_2)] = m$ and $y^T h(x_0) = 0$.□

Converse duality theorem

Again assume all assumptions of the Introduction, including (2). Denote by $F(x)$ the objective function for (P).

Suppose that (3) and (4) are satisfied for

$$(u, z, y, v, w) = (u_0, z_0, y_0, v_0, w_0) ,$$

where $y_0 \in S^*$, $v_0 \in C_1$, $w_0 \in C_2$, and $g(u_0) + w_0^T u_0 > 0$. The system

$$(3) \quad 0 \in \partial(-f+zg)(u) + \partial(y^T h)(u) + (v+zw) ,$$

together with

$$(7) \quad z = [f(u)-v^T u-y^T h(u)]/[g(u)+w^T u] ,$$

will be called *solvable near* u_0 if, whenever $y = y_0 + \beta \tilde{y} \in S^*$, $v = v_0 + \beta \tilde{v} \in C_1$, and $w = w_0 + \beta \tilde{w} \in C_2$, for $0 \leq \beta \leq 1$, then the system (3) and (7) has a solution $u = u_0 + \tilde{u}(\beta)$ for all sufficiently small positive β , satisfying $\tilde{u}(\beta) \rightarrow 0$ as $\beta \downarrow 0$.

This property holds, in particular, if $-f + z_0 g + y_0^T h$ has a non-singular hessian matrix, at $u = u_0$, in consequence of the implicit function theorem. However, the following converse duality theorem does not assume any differentiability of the functions f, g, h .

THEOREM 2. *Let (D) reach a minimum at*

$$(u, z, y, v, w) = (u_0, z_0, y_0, v_0, w_0) ,$$

where $u_0 \in X_0$. If $z_0 = 0$, assume that $F(\hat{x}) \geq 0$ for some $\hat{x} \in X_0$ satisfying $-h(\hat{x}) \in S$. If $z_0 > 0$, assume that the system (3) and (7) is solvable near u_0 . Then (P) reaches a maximum, and $\max(P) = \min(D)$. Hence (P) is a strong dual to (D).

Proof. If $z_0 = 0$, then $F(\hat{x}) \leq z_0 = 0$ by weak duality, also $F(\hat{x}) \geq 0$ by assumption. Hence $F(\hat{x}) = z_0$, and weak duality implies that \hat{x} is optimal for (P), so that $\max(P) = \min(D)$. (Weak duality is available from Theorem 1.)

Suppose now that $z_0 \neq 0$; since $z_0 \geq 0$ for (D), $z_0 > 0$. Choose any $\tilde{y}, \tilde{v}, \tilde{w}$ so that $y_0 + \tilde{y} \in S^*$, $v_0 + \tilde{v} \in C_1$, $w_0 + \tilde{w} \in C_2$. Since S^*, C_1 , and C_2 are convex sets, $y = y_0 + \beta\tilde{y} \in S^*$, $v = v_0 + \beta\tilde{v} \in C_1$, and $w = w_0 + \beta\tilde{w} \in C_2$, whenever $0 \leq \beta \leq 1$. By assumption $g(u_0) + w_0^T u_0 > 0$, and then (3) and (7) have a solution $u = u_0 + \tilde{u}(\beta)$ for sufficiently small $\beta > 0$. By continuity, $g(u) + w^T u > 0$ for sufficiently small $\beta > 0$; hence (7) implies (4). Hence this point (u, z, y, v, w) , with $u = u_0 + \tilde{u}(\beta)$ and z given by (7), is feasible for (D), for sufficiently small $\beta > 0$.

Since $(u_0, z_0, y_0, v_0, w_0)$ minimizes (D),

$$(8) \quad z_0 \leq [f(u) - v^T u - y^T h(u)] / [g(u) + w^T u] \equiv p/q,$$

using (7), where

$$(9) \quad p = f(u_0) + \beta f'(u_0; \tilde{u}) - v_0^T u_0 - \beta v_0^T \tilde{u} - \beta \tilde{v}^T u_0 - y_0^T h(u_0) - \beta y_0^T h'(u_0; \tilde{u}) - \beta \tilde{y}^T h(u_0) + o(\beta),$$

$$(10) \quad q = g(u_0) + \beta g'(u_0; \tilde{u}) + w_0^T u_0 + \beta w_0^T \tilde{u} + \beta \tilde{w}^T u_0 + o(\beta).$$

Combining these terms shows that

$$(11) \quad (p_0 - z_0 q_0) + \beta R - \beta \left[\tilde{y}^T h(u_0) + (\tilde{v} + z_0 \tilde{w})^T u_0 \right] + o(\beta) \geq 0,$$

where $p_0 = f(u_0) + y_0^T h(u_0) + v_0^T u_0$, $q_0 = g(u_0) + w_0^T u_0$, and

$$(12) \quad R = f'(u_0; \tilde{u}) - z_0 g'(u_0; \tilde{u}) - \left(y_0^T h \right)'(u_0; \tilde{u}) - v_0^T \tilde{u} - z_0 w_0^T \tilde{u}.$$

Then $p_0 - z_0 q_0 \geq 0$; but $p_0 - z_0 q_0 \leq 0$ by (4), so $p_0 - z_0 q_0 = 0$.

Dividing (11) by β and letting $\beta \rightarrow 0$, then shows that

$$(13) \quad R - \tilde{y}^T h(u_0) - (\tilde{v} + z_0 \tilde{w})^T u_0 \geq 0.$$

Denote $\psi = -f + z_0 g + y^T h$. From (3), $\theta + v_0 + z_0 w_0 = 0$ for some $\theta \in \partial\psi(u_0)$. Then $\psi'(u_0; \tilde{u}) \geq \theta^T \tilde{u}$. Since $R = \psi'(u_0; \tilde{u}) - (v_0 + z_0 w_0)^T \tilde{u}$, it follows that $R \leq -(\theta + v_0 + z_0 w_0)^T \tilde{u} = 0$. Hence

$$(14) \quad \tilde{y}^T h(u_0) + (\tilde{v} + z_0 \tilde{w})^T u_0 \leq 0.$$

Setting $\tilde{y} = 0$ and $\tilde{w} = 0$, $\tilde{v}^T u_0 \leq 0$ whenever $v_0 + \tilde{v} \in C_1$. Hence $v^T u_0 \leq v_0^T u_0$ for each $v \in C_1$. Therefore $s(u_0 | C_1) \leq v_0^T u_0$. Since $v_0 \in C_1$, by a constraint of (D), $s(u_0 | C_1) \geq v_0^T u_0$. Hence $v_0^T u_0 = s(u_0 | C_1)$. Since $z_0 > 0$, a similar argument applied to $\tilde{w}^T u_0$ shows that $w_0^T u_0 = s(u_0 | C_2)$. Now let $\tilde{v} = 0$ and $\tilde{w} = 0$. Then $\tilde{y}^T h(u_0) \leq 0$ whenever $\tilde{y} \in S^*$; since S is a closed convex cone, it follows that $-h(u_0) \in S$. Setting $\tilde{y} = -\lambda y_0$, $y_0 + \tilde{y} \in S^*$, and then $(-\lambda y_0)^T h(u_0) \leq 0$. But also $y_0^T h(u_0) \leq 0$ since $y_0 \in S^*$, and $-h(u_0) \in S$ has just been proved. Therefore $y_0^T h(u_0) = 0$.

Thus u_0 is feasible for (P), and the optimal objective function for (D) equals

$$(15) \quad z_0 = [f(u_0) + s(u_0 | C_1) - 0] / [g(u_0) + s(u_0 | C_2)] = F(u_0).$$

Using weak duality, it follows that u_0 is optimal for (P). Hence

$\max(P) = \min(D)$. \square

Discussion and examples

If f, g , and h are differentiable functions, then (3) reduces to

(16) $0 = (-f + zg + y^T h)'(u) + (v + zw)$.

For nondifferentiable functions, an equivalent to (3) is (see [15])

(17) (for all t) $\varphi'(u; t) + (v + zw)t \geq 0$,

where $\varphi = -f + zg + y^T h$.

The technique of proof for Theorem 2 is adapted from that of [9], Theorem 6, and [5], Theorem 4.8.1. Since the "solvable near u_0 " requirement in Theorem 2 does not demand a *unique* solution for u , the usual implicit function theorem is assuming too much. A general verifiable solvability criterion for the convex nondifferentiable case has yet to be found. An inclusion of the form $\rho \in \partial\Phi(u)$ must be solved (nonuniquely) for u , given $\rho_0 \in \partial\Phi(u_0)$. For this it suffices if $\partial\Phi(\cdot)$ maps a neighbourhood of u_0 onto some neighbourhood of ρ_0 . As a simple example let $\Phi(u) = (u^T u)^{\frac{1}{2}} + \frac{1}{2}\epsilon u^T u$, for $u \in \mathbf{R}^n$ and ϵ a positive constant; set $u_0 = 0$. Denote $B(r) = \{\xi \in \mathbf{R}^n : \|\xi\| \leq r\}$. Then $\partial\Phi(u_0) = B(1)$, and $\partial\Phi(\cdot)$ maps $B(\delta)$ onto $B(1 + \epsilon\delta)$, so that the sufficient requirement is fulfilled, provided that $\|\rho_0\| \leq 1$, for this nondifferentiable function Φ . This is not so if $\epsilon = 0$.

Problem (P) includes various special cases. If f, g , and h are affine, then (P) reduces to that considered by Mond and Schechter [13], [14]. As noted in [13], if B is a positive semidefinite matrix, and Q is the compact set $\{Bv : v^T Bv \leq 1\}$, then

(18) $s(x|Q) = (x^T Bx)^{\frac{1}{2}}$.

Thus, if f, g , and h are differentiable functions, and $s(\cdot|C_1)$ and $s(\cdot|C_2)$ are defined as in (18) by positive semidefinite matrices B and D , problem (P) becomes the nondifferentiable fractional programming

problem considered in [11]. If also f, g , and h are affine, the results of Chandra and Gulati [3] are obtained. If $g(x) \equiv 1$, and $C_2 = \{0\}$, then $s(x|C_2) = 0$, so that a nonfractional nondifferentiable objective function $F(x) = f(x) - s(x|C_1)$ is recovered.

If S is a $k \times n$ matrix, then [13]

$$(19) \quad \|Sx\|_p = s(x|Q)$$

for $p \geq 1$, where $Q = \{S^T u : \|u\|_q \leq 1\}$, and $p^{-1} + q^{-1} = 1$, and

$q = \infty$ if $p = 1$. Here $\|x\|_p = \left[\sum |x_i|^p \right]^{1/p}$ if $p < \infty$, and $\|x\|_\infty = \sup\{|x_i| : i = 1, 2, \dots\}$. The set Q , as defined, is convex and compact.

Thus if f, g , and h are differentiable, $s(\cdot|C_1)$ is defined as in (19) by a matrix S_1 and a scalar p_1 , and similarly $s(\cdot|C_2)$ by a matrix S_2 and a scalar p_2 , then problems (P) and (D) become respectively

$$(P'): \text{ maximize } [f(x) - \|S_1 x\|_{p_1}] / [g(x) + \|S_2 x\|_{p_2}] \text{ subject to } h(x) \leq 0, \\ x \in X_0$$

$$(D'): \text{ minimize } z \text{ subject to } z \geq 0, y \geq 0, \|v\|_{q_1} \leq 1, \|w\|_{q_2} \leq 1, \\ u \in X_0, z, y, v, w$$

$$\nabla(y^T h - f + z g)(u) + S_1^T v + S_2^T w = 0,$$

$$-f(u) + z g(u) + y^T h(u) + u^T \left\{ S_1^T v + S_2^T w \right\} \geq 0.$$

If, in particular, f, g , and h are affine functions, then some of the problems discussed in [13] are obtained.

If f, g , and h are differentiable, and C_2 consists only of the zero vector in \mathbf{R}^n , and $s(\cdot|C_1)$ is defined, as in (18), by a positive semidefinite matrix B , then (P) and (D) reduce to the problems considered by Aggarwal and Saxena [1], [2]. If also $g(x) \equiv 1$, one obtains the (nonfractional) problems discussed in [10].

If f, g , and h are differentiable, $C_2 = \{0\}$, $g(x) \equiv 1$, and $s(\cdot|C_1)$ is defined as in (19) by a matrix S , then the present results yield those of Mond and Schechter [12].

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