

# THE DENSEST PACKING OF SIX SPHERES IN A CUBE

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This packing problem is obviously equivalent to the problem of locating six points  $P_i (1 \leq i \leq 6)$  in a closed unit cube  $C$  such that  $\min_{i \neq j} d(P_i, P_j)$  is as large as possible, where  $d(P_i, P_j)$  denotes the distance between  $P_i$  and  $P_j$ . We shall prove that this minimum distance cannot exceed  $\frac{3\sqrt{2}}{4}$  ( $=m$ , say), and that it attains this value only if the points form a configuration which is congruent to the one of the points  $R_i (1 \leq i \leq 6)$  shown in fig. 1. Note that  $d(R_i, A_i) = \frac{1}{4} (1 \leq i \leq 6)$ , and so the six points are the vertices of a regular octahedron.

1) For our proof we shall need the solution of the analogous problem for three points in a right square prism  $P$  of side 1 and height  $\frac{1}{4}$ :  $0 \leq y_i \leq 1 (i=1, 2), 0 \leq y_3 \leq \frac{1}{4}$ .

PROPOSITION 1. For any three points  $Q_1, Q_2, Q_3$  of  $P$ ,  $\min_{i \neq j} d(Q_i, Q_j) \leq \frac{3\sqrt{2}}{4} = m$ , and equality holds only for a configuration congruent to the set of the points  $V_1(\frac{1}{4}, 1, \frac{1}{4})$ ,  $V_2(1, \frac{1}{4}, \frac{1}{4})$ , and  $V_3(0, 0, 0)$ . See fig. 2. Note that  $d(V_i, V_j) = m (i \neq j)$ .

Proof. Consider any best configuration<sup>1</sup>  $T$  of three points  $Q_1, Q_2, Q_3$  in  $P$ . Of course

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<sup>1</sup> i. e., a configuration for which  $\min_{i \neq j} d(Q_i, Q_j)$  is maximum.

$$(1) \quad \min_{i \neq j} d(Q_i, Q_j) \geq m.$$

(A) Assume first that a point of  $T$  lies in a vertex of  $P$ , say  $Q_3 = V_3$ . Then by (1) no other point of  $T$  can lie in the convex hull  $H$  of the vertices  $V_3, (0, 0, \frac{1}{4}), (1, 0, 0), (1, 0, \frac{1}{4}), (0, 1, 0), (0, 1, \frac{1}{4})$  of  $T$ ,  $V_1, V_2, U_1(\frac{\sqrt{2}}{4}, 1, 0)$ , and  $U_2(1, \frac{\sqrt{2}}{4}, 0)$ , except possibly at  $V_1, V_2, U_1$ , or  $U_2$ . Note that  $d(V_3, V_i) = d(V_3, U_i) = m (i=1, 2)$ . Therefore  $Q_1$  and  $Q_2$  must lie in the closure of  $P - H$ . But this polyhedron<sup>2</sup> assumes its diameter  $m$  only between the points  $V_1$  and  $V_2$ . Therefore  $\{Q_1, Q_2\} = \{V_1, V_2\}$ .

(B) We are left to show that at least one point of  $T$  must lie in a vertex of  $P$ . If we assume the contrary, then  $Q_1, Q_2$ , and  $Q_3$  must lie on mutually orthogonal non-intersecting edges of  $P$ . This follows from the basic lemma according to which on every face of  $P$  there must be at least one point of any best configuration [1]. Thus we may assume  $Q_1 = (y_1, 1, \frac{1}{4})$ ,  $Q_2 = (1, y_2, 0)$ , and  $Q_3 = (0, 0, y_3)$ , with  $0 < y_i < 1 (i=1, 2)$ , and  $0 < y_3 < \frac{1}{4}$ . By (1)  $d^2(Q_3, Q_i) > m^2 (i=1, 2)$ . This leads to

$$y_1 > \sqrt{\frac{1}{8} - (\frac{1}{4} - y_3)^2} \quad \text{and} \quad y_2 > \sqrt{\frac{1}{8} - y_3^2}.$$

$$\begin{aligned} \text{But then } d^2(Q_1, Q_2) &= (1-y_1)^2 + (1-y_2)^2 + \frac{1}{16} \\ &< 2 + \frac{1}{2}y_3 - 2y_3^2 - 2\sqrt{\frac{1}{8} - (\frac{1}{4} - y_3)^2} - 2\sqrt{\frac{1}{8} - y_3^2}. \end{aligned}$$

For  $0 < y_3 < \frac{1}{4}$  this expression is less than  $\frac{9}{8}$ , in contradiction to (1), q. e. d.

This proves that if three points with mutual distances at least  $m = \frac{3\sqrt{2}}{4}$  lie in a right square prism of side 1, then the height of the prism must be at least  $\frac{1}{4}$ .

<sup>2</sup> The diameter of a closed polyhedron is obviously always assumed between two of its vertices.

2) Let  $S$  be any set of six points  $P_i (1 \leq i \leq 6)$  of  $C$  such that

$$(2) \quad d(P_i, P_j) \geq m \quad (1 \leq i < j \leq 6)$$

We shall prove that  $\{R_i (1 \leq i \leq 6)\}$  of fig. 1 is, up to congruent ones, the only such set.

The unit cube  $C: 0 \leq x_j \leq 1 (1 \leq j \leq 3)$  is the union of eight closed cubes  $C_k$  of side  $\frac{1}{2}$ :  $a_j \leq x_j \leq b_j$ , where either  $a_j = 0$  and  $b_j = \frac{1}{2}$ , or  $a_j = \frac{1}{2}$  and  $b_j = 1$ . Let us enumerate them such that the vertices  $A_k \in C_k (1 \leq k \leq 8)$  (see fig. 1). Since  $\frac{\sqrt{3}}{2} < m$ , by (2) in every  $C_k$  there can at most be one point of  $S$ . Therefore we may choose six cubes  $C_k$  containing one point of  $S$  each, and two "empty" cubes that do not contain any point of  $S$  except possibly on their intersection with a "containing" cube. This choice may be not unique, but all we need is the existence of (at least) two such "empty" cubes  $C_k$ .

**PROPOSITION 2.** If two "containing" cubes  $C_i, C_j$  are adjacent, then the points of  $S$  which they contain have at least a distance  $\frac{1}{2} - a$  from their common face,  $a \equiv 1 - \frac{\sqrt{10}}{4} < \frac{1}{4}$ .

Indeed, consider the right square prism of side  $\frac{1}{2}$  and diagonal  $m$  which contains all of  $C_i$  and as much of  $C_j$  as possible (see fig. 3). Excluding its base, which lies completely in  $C_j$ , it can, because of (2), contain at most one point of  $S$ . But it contains already the point of  $S$  in  $C_i$ . Therefore the point of  $S$  in  $C_j$  must lie in the indicated right square prism of side  $\frac{1}{2}$  and height  $a$ ,  $a$  being defined by  $(1-a)^2 + 2(\frac{1}{2})^2 = m^2$ .  
q. e. d.

**COROLLARY.** If a "containing" cube  $C_k$  is adjacent to two other "containing" cubes, then the point of  $S$  in  $C_k$  is confined in a right square prism of side  $a$  and height  $\frac{1}{2}$  with one edge common with that edge of  $C_k$  which has no points in

common with the two adjacent "containing" cubes.

3) PROPOSITION 3. The six points  $P_i$  must lie in right square prisms of side  $a (< \frac{1}{4})$  and height  $\frac{1}{2}$ , namely (see fig. 4)  $P_i$  in  $\frac{1}{2} \leq x_i \leq 1$ ,  $0 \leq x_j \leq a$  ( $j \neq i$ ) ( $i=1, 2, 3$ ) resp. in  $0 \leq x_{i-3} \leq \frac{1}{2}$ ,  $1 - a \leq x_j \leq 1$  ( $j \neq i-3$ ) ( $i=4, 5, 6$ )

Proof. The two "empty" cubes  $C_k$  are not adjacent, nor can they have a common edge. Otherwise their centers would have at least one equal coordinate, say  $x_3 = \frac{3}{4}$ , and the four cubes  $C_k: 0 \leq x_3 \leq \frac{1}{2}$  would all be "containing". By the corollary of Proposition 2 the points of  $S$  which they contain would be confined to  $0 \leq x_j \leq a$  or  $1 - a \leq x_j \leq 1$  ( $j=1, 2$ ). The four right square prisms of side  $a$  and height  $\frac{1}{4}: 0 \leq x_j \leq \frac{1}{4}$  or  $\frac{3}{4} \leq x_j \leq 1$  ( $j=1$  or  $2$ ),  $0 \leq x_h \leq 1$  ( $h \neq j$ ), would therefore already contain at least 2 points of  $S$  each. Thus by Proposition 1 the two other points of  $S$ , i.e. those with  $\frac{1}{2} \leq x_3 \leq 1$ , would be restricted to  $\frac{1}{4} \leq x_j \leq \frac{3}{4}$  ( $j=1, 2$ ). But this is impossible, because this point set is a cube of side  $\frac{1}{2}$  and diameter  $\frac{\sqrt{3}}{2} < m$ .

Thus the two "empty" cubes must lie opposite to the center of  $C$ , e.g. let them be  $C_7$  and  $C_8$ . We may then assume  $P_i \in C_i$  ( $1 \leq i \leq 6$ ). Proposition 3 follows now from the corollary of Proposition 2.

4) Using Proposition 3 and applying Proposition 1 to the six right square prisms in  $C$  of height  $\frac{1}{4}$ , which contain one face of  $C$  each, the location of the  $P_i$  can in addition be restricted to  $P_1, P_2, P_3$  all in  $0 \leq x_j \leq \frac{3}{4}$  ( $1 \leq j \leq 3$ ),  $P_4, P_5, P_6$  all in  $\frac{1}{4} \leq x_j \leq 1$  ( $1 \leq j \leq 3$ ).

5) According to the solution of the analogous problem of placing three points in a cube [1] the only way to locate three

points with minimum distance  $\frac{3\sqrt{2}}{4}$  in a cube of side  $\frac{3}{4}$  consists

in placing them in vertices with mutual distances  $\frac{3\sqrt{2}}{4}$ .

Applying this result to 4) and Proposition 3 we deduce  $P_i = R_i$  ( $1 \leq i \leq 6$ ).

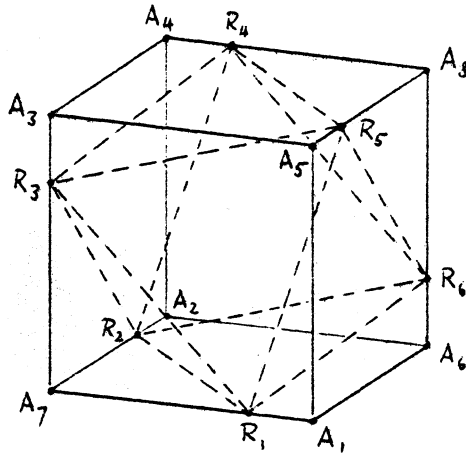


Figure 1

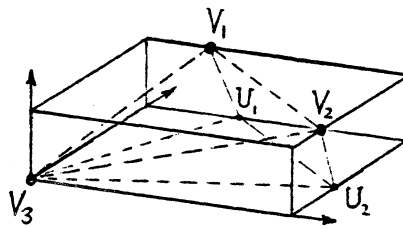


Figure 2

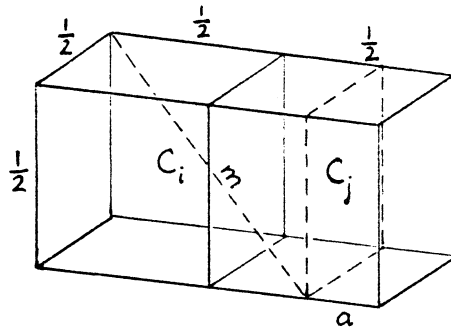


Figure 3

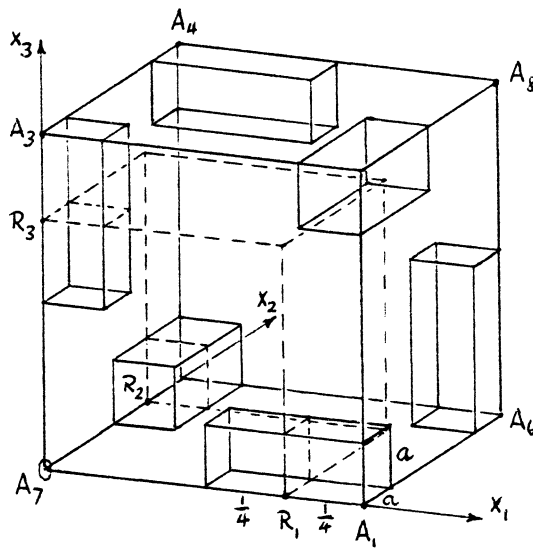


Figure 4

#### REFERENCES

1. J. Schaer, On the densest packing of spheres into a cube. *Canad. Math. Bull.* vol. 9, no. 3, 1966.

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